Homework 1

Due: Thursday, Jan 30, 2025 at 12:00pm (Noon)

Written Assignment

Introduction

The purpose of this assignment is to fortify your background in statistics, linear algebra, runtime complexity, socially responsible computing, and programming with NumPy. A strong grasp of these subjects is essential for your success in the course.

You may be able to find answers to these problems by searching for the problem text or consulting a large language model. Please search instead for the concepts being applied; the goal is not to solve these specific problems but to review principles that will be fundamental to the course.

Problem 0: Collaboration Policy

Every student enrolled in CSCI 1420 is expected to know and abide by the course collaboration policy. Your assignments will not be graded if you do not meet this expectation.

Please review the course collaboration policy available at this link and include the following signed statement as your answer to Problem 0:

By signing with my Banner ID below, I agree to abide by the CSCI1420 Collaboration Policy as published on the course website. I also agree to not include my real name on assignments in the course and instead will use my Banner ID on any applicable submissions to preserve anonymity in grading.

Signed: *Your Banner ID Here*

Problem 1: Review of Probability and Statistics

(12 points)

a. Recall Bayes' Theorem, which states that

$$P(B \mid A) = \frac{P(A \mid B) \cdot P(B)}{P(A)} = \frac{P(A \cap B)}{P(A)}.$$

for a probability distribution P concerning events A and B.

Brown students love pastries. Suppose we have a distribution P that describes the characteristics of Brown students' pastries, where P(pastry is cake) = P(pastry is chocolate-flavored) = 1/2. Additionally, P(pastry is expensive) = 1/7. Assume that each Brown student with a pastry has exactly one pastry. Assume that encounters with Brown students with pastries are independent events unless otherwise specified.

i. Steve bumps into two random Brown students with pastries. What is the probability that both of their pastries are cakes?

Solution:

$$P(\text{both cakes}) = P(\text{first is cake} \cap \text{second is cake})$$
$$= P(\text{first is cake}) \cdot P(\text{second is cake})$$
$$= \frac{1}{2} \cdot \frac{1}{2} = \boxed{\frac{1}{4}}$$

ii. Jania bumps into two random Brown students with pastries. One is a computer science concentrator, and the other is a history concentrator. Suppose the history concentrator's pastry is a cake. What is the probability that both of their pastries are cakes?

Solution:

$$P \text{ (both cakes | first is cake)} = \frac{P(\text{both cakes}) \cap P(\text{first is cake})}{P(\text{first is cake})}$$
$$= \frac{1/4}{1/2} = \boxed{\frac{1}{2}}$$

iii. Prithvi bumps into two random Brown students with pastries. At least one of the students' pastries is chocolate-flavored. What is the probability that both of their pastries are chocolate-flavored?

Solution:

 $P \text{ (both chocolate-flavored | at least one is chocolate-flavored)} = \frac{P \text{ (both chocolate-flavored)} \cap P \text{ (at least one is chocolate-flavored)}}{P \text{ (at least one is chocolate-flavored)}} = \frac{1/4}{3/4} = \boxed{\frac{1}{3}}$

iv. Andy bumps into two random Brown students with pastries. At least one of the students' pastries is an expensive cake. What is the probability that both of their pastries are cakes?

Solution: Let's say that the first pastry is pastry A and the second pastry is pastry B. Let H_A denote the event that A is a cake, and let L_A denote the event that A is expensive. Similarly, let H_B denote the event that B is a cake, and let L_B denote the event that B is expensive. We're given that at least one pastry is a cake and is expensive. We can represent this as the following event: $(H_A \cap L_A) \cup (H_B \cap L_B)$. We wish to find the probability that both pastries are cakes, or $P(H_A \cap H_B | (H_A \cap L_A) \cup (H_B \cap L_B))$.

$$P\left(H_A \cap H_B \mid (H_A \cap L_A) \cup (H_B \cap L_B)\right) = \frac{P\left((H_A \cap H_B) \cap ((H_A \cap L_A) \cup (H_B \cap L_B))\right)}{(H_A \cap L_A) \cup (H_B \cap L_B)}$$

Let's expand the event in the numerator.

$$(H_A \cap H_B) \cap ((H_A \cap L_A) \cup (H_B \cap L_B))$$

= $((H_A \cap H_B) \cap (H_A \cap L_A)) \cup ((H_A \cap H_B) \cap (H_B \cap L_B))$
= $(H_A \cap H_B \cap L_A) \cup (H_A \cap H_B \cap L_B)$
= $(H_A \cap H_B) \cap (L_A \cup L_B)$

Notice that the event $L_A \cup L_B$ occurs when any one of the events $L_A \cap L_B^c$, $L_A^c \cap L_B$, or $L_A \cap L_B$ occur. These events are mutually exclusive. So,

$$P(L_A \cup L_B) = P(L_A \cap L_B^c) + P(L_A^c \cap L_B) + P(L_A \cap L_B)$$

= $\frac{1}{7} \cdot \frac{6}{7} + \frac{6}{7} \cdot \frac{1}{7} + \frac{1}{7} \cdot \frac{1}{7}$
= $\frac{13}{49}$

We can then find the value of the numerator.

$$P\left((H_A \cap H_B) \cap \left((H_A \cap L_A) \cup (H_B \cap L_B)\right)\right) = P\left((H_A \cap H_B) \cap (L_A \cup L_B)\right)$$
$$= P\left(H_A \cap H_B\right) P\left(L_A \cup L_B\right)$$
$$= \frac{1}{4} \cdot \frac{13}{49}$$
$$= \frac{13}{196}$$

Now, on to the denominator.

$$P\left((H_A \cap L_A) \cup (H_B \cap L_B)\right) = 1 - P\left((H_A \cap L_A)^c \cap (H_B \cap L_B)^c\right)$$

= 1 - P ((H_A \cap L_A)^c) P ((H_B \cap L_B)^c)
= 1 - (1 - P (H_A \cap L_A))(1 - P (H_B \cap L_B))
= 1 - \left(1 - \frac{1}{2} \cdot \frac{1}{7}\right) \left(1 - \frac{1}{2} \cdot \frac{1}{7}\right)
= 1 - $\frac{169}{196}$
= $\frac{27}{196}$

Putting it all together, we end up with

$$P(H_A \cap H_B | (H_A \cap L_A) \cup (H_B \cap L_B)) = \frac{P((H_A \cap H_B) \cap ((H_A \cap L_A) \cup (H_B \cap L_B)))}{(H_A \cap L_A) \cup (H_B \cap L_B)}$$
$$= \frac{P((H_A \cap H_B) \cap (L_A \cup L_B))}{P((H_A \cap L_A) \cup (H_B \cap L_B))}$$
$$= \frac{13/196}{27/196}$$
$$= \boxed{\frac{13}{27}}$$

Hint: None of the four parts have the same answer.

b. Suppose Z_1, \ldots, Z_k are independent standard normal random variables. This means that $\mathbb{E}[Z_i] = 0$ and $\operatorname{Var}(Z_i) = 1$. Let

$$Q = \sum_{i=1}^{k} Z_i^2.$$

Derive an expression for $\mathbb{E}[Q]$ in terms of k.

Hint: Apply the linearity of expectation.

Remark: Q as defined above follows the Chi-Squared Distribution.

Solution: By the linearity of expectation,

$$\mathbb{E}[Q] = \mathbb{E}\left[\sum_{i=1}^{k} Z_i^2\right]$$
$$= \sum_{i=1}^{k} \mathbb{E}\left[Z_i^2\right].$$

Recall that the variance of a random variable Z may be expressed as

$$\operatorname{Var}(Z) = \mathbb{E}\left[Z^2\right] - \left(\mathbb{E}[Z]\right)^2.$$

Since Z_j is standard normal, $\mathbb{E}[Z_j] = 0$, and $\operatorname{Var}(Z_j) = 1$. Therefore, $\mathbb{E}[Z_j^2] = 1$. Plugging this result in, we find that

$$\mathbb{E}[Q] = \sum_{i=1}^{k} 1$$
$$= \boxed{k}.$$

Problem 2: Review of Linear Algebra

(12 points)

a. Consider a matrix $A \in \mathbb{R}^{n \times n}$. Let rank(A) < n. Prove that A has an eigenvalue equal to 0. You may not use the Invertible Matrix Theorem when solving this problem.

Hint: Recall that the rank of a matrix is equal to the number of linearly independent columns of the matrix.

Solution 1 (Linear Dependence): Let $\mathbf{A}_1, ..., \mathbf{A}_n$ be the columns of A. Since rank(A) < n, the n columns of A are linearly dependent. Therefore there is linear combination of the columns of A that is equal to the zero vector. In other words, there exist $k_1, ..., k_n$ such that $k_1\mathbf{A}_1 + k_2\mathbf{A}_2 + ... + k_n\mathbf{A}_n = \mathbf{0}$. Then, $A\mathbf{k} = \mathbf{0}$, where

$$\mathbf{k} = \begin{bmatrix} k_1 \\ \dots \\ k_n \end{bmatrix}$$

We can then conclude that $A\mathbf{k} = 0 \cdot \mathbf{k}$, and 0 is an eigenvalue of A by definition.

Solution 2 (Rank-Nullity): We use the Rank-Nullity Theorem, which gives

$$\operatorname{rank}(A) + \operatorname{nullity}(A) = \dim(\mathbb{R}^n) = n$$

Since $\operatorname{rank}(A) < n$, we have $\operatorname{nullity}(A) = n - \operatorname{rank}(A) > 0$. This means there is some non-zero vector, v, in the null space of A. By definition, we know that Av = 0 = 0v, so 0 is an eigenvalue of A (and v is an associated eigenvector).

b. Let $B \in \mathbb{R}^{n \times n}$ be a symmetric matrix satisfying $\mathbf{x}^T B \mathbf{x} > 0$ for all $\mathbf{x} \in \mathbb{R}^n$ such that $\mathbf{x} \neq \mathbf{0}$. Remember that the eigenvalues of a symmetric matrix are real. Prove that all of the eigenvalues of B are strictly greater than 0. Do not use any theorems for positive definite matrices.

Solution: Let λ be an eigenvalue of B and let **v** be the corresponding eigenvector. Then,

$$B\mathbf{v} = \lambda \mathbf{v}$$

Next, we can multiply by \mathbf{v}^T on both sides:

$$\mathbf{v}^T B \mathbf{v} = \mathbf{v}^T \lambda \mathbf{v} = \lambda \mathbf{v}^T \mathbf{v} = \lambda ||\mathbf{v}||^2.$$

We know that $\mathbf{v}^T B \mathbf{v}$ is strictly greater than 0. The square of the norm of \mathbf{v} is also strictly greater than 0, because eigenvectors are nonzero. Therefore, λ must also be greater than 0. Thus, all of the eigenvalues of B are greater than 0.

c. Let $C \in \mathbb{R}^{n \times n}$ be a symmetric matrix. Assume that C has n unique eigenvalues. Prove that all of the eigenvectors of C are orthogonal.

Hint: Begin your proof by considering the equations

$$C\mathbf{v}_1 = \lambda_1 \mathbf{v}_1$$
$$C\mathbf{v}_2 = \lambda_2 \mathbf{v}_2$$

where $(\lambda_1, \mathbf{v}_1)$ and $(\lambda_2, \mathbf{v}_2)$ are two different eigenpairs of *C*. Then, multiply the second equation by \mathbf{v}_1^T . Solution: Let λ_1 and λ_2 be different eigenvalues of *C* and let \mathbf{v}_1 and \mathbf{v}_2 be the corresponding eigenvectors. So,

$$C\mathbf{v}_1 = \lambda_1 \mathbf{v}_1$$
 and $C\mathbf{v}_2 = \lambda_2 \mathbf{v}_2$

Then, multiply the second equation by \mathbf{v}_1^T :

$$\mathbf{v}_1^T C \mathbf{v}_2 = \lambda_2 \mathbf{v}_1^T \mathbf{v}_2$$
$$(C^T \mathbf{v}_1)^T \mathbf{v}_2 = \lambda_2 \mathbf{v}_1^T \mathbf{v}_2$$
$$(C \mathbf{v}_1)^T \mathbf{v}_2 = \lambda_2 \mathbf{v}_1^T \mathbf{v}_2$$
$$\lambda_1 \mathbf{v}_1^T \mathbf{v}_2 = \lambda_2 \mathbf{v}_1^T \mathbf{v}_2$$
$$(\lambda_1 - \lambda_2) \mathbf{v}_1^T \mathbf{v}_2 = 0$$

We know that $\lambda_1 - \lambda_2 \neq 0$ because $\lambda_1 \neq \lambda_2$. So, $\mathbf{v}_1^T \mathbf{v}_2 = 0$. Therefore, all of the eigenvectors of C are orthogonal.

Problem 3: Runtime Complexity

(14 points)

The Fibonacci sequence can be defined by the recurrence relation F(n) = F(n-1) + F(n-2) for $n \ge 2$, with F(0) = 0 and F(1) = 1. We can make use of the fact (likely first noted by Edsger Dijkstra) that

$$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^n = \begin{pmatrix} F(n+1) & F(n) \\ F(n) & F(n-1) \end{pmatrix}$$

for any positive integer *n*. So, to compute F(n), we can compute $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^{n-1}$ and return the upper left number in the matrix.

Describe an algorithm to compute F(n) for arbitrary n that uses <u>fewer than</u> O(n) additions and multiplications (don't worry about the size of the integers), and prove that it satisfies this complexity bound. That is, prove a sublinear complexity bound in terms of n on the number of additions and multiplications this algorithm performs. Feel free to search for such an algorithm, but please generate the proof of the complexity bound yourself.

To compute F(n) we can compute $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^{n-1}$ and return the upper left number. To compute $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^n$ for even n, we just need to compute $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^{\frac{n}{2}}$ and multiply it by itself - in other words, we can reduce the exponent by half with a constant amount of work. This suggests an algorithm with complexity $O(\log n)$.

Pseudo-code:

function $FIB(n)$
if $n = 0$ then
return 0
else if $n = 1$ then
return 1
$mat \leftarrow POWER(n-1)$
return mat[0][0]

1: function $POWER(n)$		
2:	if $n = 1$ then	
3:	return $[[1,1],[1,0]]$	
4:	else if n is even then	
5:	$mat \leftarrow POWER(n/2)$	
6:	$_$ return mat * mat	
7:	$mat \leftarrow POWER(n-1)$	
8:	return <i>mat</i> * [[1,1],[1,0]]	

Proof of time complexity:

A call to fib makes at most one call to power. Each call to power results in at most one recursive call. The total number of recursive calls is $O(\log n)$ because we divide n by 2 at least every other call (in the worst case, we get an odd number every time we divide by 2, but then we immediately make a recursive call with an even number and divide by 2). We do at most a constant number of arithmetic operations in each recursive call of power, and at most constant operations in fib, so the algorithm has $O(\log n)$ arithmetic operations.

Problem 4: Socially Responsible Computing

(7 points)

While planning the construction of the Fukushima nuclear power plant, data scientists were interested in determining the highest earthquake magnitude the plant should be able to withstand. To achieve this, they

used continuous regression models using data of earthquake magnitudes recorded over the years. The graphs below plot two models, where the y-axis represents annual frequency and the x-axis represents earthquake magnitude. The magnitude ranges 4.5 up to the rare 9.5. The diamonds denote the observed frequencies of earthquakes of different magnitudes, and the solid line represents the predictions of the frequencies by each model.



Source: Silver, N. (2012). The Signal and the Noise: Why So Many Predictions Fail-but Some Don't. New York, NY: Penguin Press.

Scientists referred to one of the two graphs above to inform the plant's engineering and determined that the plant should withstand an earthquake of up to 8.6. In March 2011, the Fukushima nuclear power plant was hit with a 9.1 earthquake—one the most severe nuclear accidents since Chernobyl.

a. The scientists' determination for the plant's design to withstand an earthquake of 8.6 magnitude led to a detrimental consequence. Looking at the two graphs, what does each model predict about the probability of an earthquake of magnitude 8.6?

Solution: This question is designed to get students thinking about how researchers were using this model to predict/understand something they have never seen before (an earthquake of greater than magnitude 9). Some key points/example explanations to look out for in answers:

- Models were both looking at frequency of earthquake vs. magnitude of earthquake.
- Left graph uses a characteristic fit, where the trend line fits the training data very well.
- Right graph uses a Gutenberg-Richter fit, where trend line does not fit as well
- Model in right graph predicts a higher frequency than model in left graph for higher magnitude earthquakes
- b. Which graph did the scientists use in making their decision for the plant's design? What went wrong in the modeling, and why should they have referred to the other graph instead?

Hint: Think of things that often go wrong with choosing and fitting a model to some given data, and make note of the difference in the two graphs.

Solution:

- The scientists used the **right graph** in making their decision for the plant's design
- What to look for in answers: **overfitting**

- The disaster was a result of dangers of overfitting used the model in the second graph to inform decisions, which gave an accurate prediction for training data but not in real world application.
- Scientists should have referred to the right side graph because model is not as fitted to data, and can generalize better

Programming Assignment

Introduction

The purpose of this portion is to get familiar with NumPy, which we will use heavily throughout the semester.

NumPy is a powerful numerical computing library for Python that provides very efficient, heavily optimized functionalities for large, multi-dimensional arrays and matrices (see np.ndarray), as well as many standard manipulations and linear algebra algorithms (see np.linalg).

It is strongly preferred over "pure Python" implementations (eg. with for-loops) for several reasons. In particular, NumPy leverages several highly optimized low-level libraries (see BLAS and LAPACK) written in languages like C and Fortran to minimize all overhead involved in the operations and computations, since Python is an interpreted language whereas C and Fortran are compiled. Additionally, NumPy arrays will always be much more memory-efficient than Python list's.

Set up

Python 3.9 and NumPy are necessary for the completion of programming assignments for this course. Python and NumPy can be installed on your own machine via the directions found at this link. For your convenience, we also have a course-wide virtual environment set up on the department machines at /course/cs1420/cs142_env. The environment can be activated from your own folder by running the following command: source /course/cs1420/cs142_env/bin/activate.

Stencil Code

You can find the stencil code for this assignment on GitHub classroom at this link. For more details, please see the download/submission guide.

We have provided the following stencil code:

- warmup.py contains 8 functions for you to implement: part_a through part_g. You will implement these functions in the Warm Up: Introduction to NumPy portion of this programming assignment.
- eigenvalues.py contains two functions for you to implement: qr and compute_eigenvalues. You will implement these functions in the Computing Eigenvalues via QR Decomposition portion of this programming assignment.

You should only need to modify code marked with **#TODO** in **warmup.py** and **eigenvalues.py** to complete this assignment. If you make any other edits to the provided stencil code, please make sure to revert them or leave them commented out in the final handin for autograder compatibility.

To test your implementations in eigenvalues.py, you can run python tests.py in a terminal. Make sure you activate the virtual environment first when working over ssh or on a department machine:

source /course/cs1420/cs142_env/bin/activate

Warm Up: Introduction to NumPy

If you have never used NumPy before, please read through the "Getting started" portion of the NumPy user guide before attempting this exercise. For each of the following parts, please implement its associated function in warmup.py.

NumPy is often abbreviated as "np" (i.e., import numpy as np). You may find the following functions particularly helpful: np.arange, np.zeros, np.ones, np.eye, np.sum, np.hstack, np.vstack, np.transpose, np.matmul, np.inner, np.where, and np.dot.

You may not use np.array() (eg. to create the desired Numpy ndarray's from Python list's).

- a. Create a 1-dimensional matrix containing the values 2 through 6 inclusive. Return the new matrix in function part_a.
- b. Create a 4×4 matrix where all entries are 1. Return the new matrix in function part_b.
- c. Create a 6×6 identity matrix. Return the new matrix in function part_c.
- d. Given matrix A, compute the sum of the values of its columns. Return the answer in function part_d.
- e. Given matrices A and B, compute C, where $C = A^T B$. Return the matrix C in function part_e.
- f. Using either np.vstack or np.hstack, create a 4×2 matrix where all the values in the first column are zeros and all the values in the second column are ones. Return the new matrix in function part_f.
- g. Given a matrix of floats A, create a new matrix of the same shape where all values greater than 3.0 are set to 1 and all other values are set to 0. Return the new matrix in function part_g.

Computing Eigenvalues via QR Decomposition

Your task is to finish implementing both the Householder QR decomposition algorithm (lines 2 through 9 in Algorithm 1) and the QR eigenvalue computation algorithm (Algorithm 2). In the stencil, these are contained respectively in qr and compute_eigenvalues in the eigenvalues.py. The lines in Algorithm 1 which you must implement are colored violet. Please note that this problem is intended to serve as a NumPy exercise; therefore, understanding the theory behind the algorithms is not a requirement.

QR Decomposition

Let $A \in \mathbb{R}^{n \times n}$. The QR decomposition of A is a matrix product QR, where $Q \in \mathbb{R}^{n \times n}$ is an orthogonal matrix $(Q^T = Q^{-1})$ and $R \in \mathbb{R}^{n \times n}$ is upper triangular. This factorization is called the QR decomposition.

Below, we present the Householder QR decomposition algorithm. Again, you do not need to understand the mathematics behind this algorithm. The purpose of implementing this algorithm is to gain familiarity with transforming mathematical and algorithmic notation into respective NumPy operations. This includes learning how to multiply matrices, compute norms, and manipulate NumPy arrays using NumPy's functionalities instead of Python for-loops.

Some notes regarding the notation presented in Algorithm 1:

- 1. We let y_i denote the *i*-th entry of some vector \mathbf{y} , and $\mathbf{y}_{a:b}$ will denote the subvector \mathbf{y} with first element y_a and last element y_{b-1} . Note that the vector $\mathbf{y}_{a:b}$ should contain b-a entries.
- 2. Similarly, $x_{i,j}$ will denote the entry in the *i*-th row and *j*-th column of some matrix X. Let $X_{a:b,c:d}$ denote the submatrix of X with upper left corner $x_{a,b}$ and lower right corner $x_{c-1,d-1}$. Note that $X_{a:b,c:d}$ should have dimension $(c-a) \times (d-b)$. You can find some examples of this notation on our NumPy guide.
- 3. Also, $X_{a:b,j}$ will denote the *j*-th column of submatrix $X_{a:b,0:n}$. Note that $X_{a:b,j}$ is a vector with b a entries.
- 4. Let $I^{n \times n}$ denote the $n \times n$ identity matrix. Let $0^{n \times n}$ denote an $n \times n$ matrix where all entries are 0.
- 5. $\|\cdot\|_2$ denotes the L2 norm. sign(x) evaluates to 1 if $x \ge 0$ and 0 otherwise.

Algorithm 1 Householder QR Decomposition Algorithm

```
1: function QR(A)
             R \leftarrow \operatorname{copy}(A)
 2:
             V \leftarrow 0^{n \times n}
 3:
             I \leftarrow I^{n \times n}
 4:
             for k = 0 to n - 1 do
 5:
                   \mathbf{x} \leftarrow R_{k:n,k}
 6:
                   V_{k:n,k} \leftarrow \operatorname{sign}(x_0) \|\mathbf{x}\|_2 I_{0:n-k,0} + \mathbf{x}
 7:
                   V_{k:n,k} \leftarrow V_{k:n,k} / \|V_{k:n,k}\|_2
 8:
                   R_{k:n,k:n} \leftarrow R_{k:n,k:n} - 2V_{k:n,k}(V_{k:n,k}^T R_{k:n,k:n})
 9:
             \bar{Q} \gets 0^{n \times n}
10:
             for j = 0 to n - 1 do
11:
12:
                   \mathbf{x} \leftarrow I_{0:n,i}
                   for k = n - 1 downto 0 do
13:
                         \mathbf{x}_{k:n} \leftarrow \mathbf{x}_{k:n} - 2V_{k:n,k}(V_{k:n,k}^T \mathbf{x}_{k:n})
14:
                   \overline{Q}_{0:n,j} \leftarrow \mathbf{x}
15:
             return Q, R
16:
```

Some notes regarding the implementation of Algorithm 1:

- 1. The sign(\cdot) function has been implemented for you in the stencil code. Do not use np.sign.
- The NumPy equivalent for X_{a:b,c:d} is X[a:b, c:d] (see Python's *slice notation*). Similarly, the NumPy equivalent for X_{a:b,c} is X[a:b, c]. If you wish to index the entire j-th column of X, use X[:, j].
- 3. The vector V_{k:n,k} in line 9 should be treated strictly as a column vector. Note that calling V[k:n, k] will return a 1-dimensional NumPy array. To transform V_{k:n,k} into a column vector, you can call V[k:n, k].reshape(-1, 1). Generally, if x satisfies x.shape = (a,), then you can reshape x into an a × 1 NumPy array by calling x.reshape(-1, 1).shape = (a, 1).
- 4. You should not use V[k:n, k] /= np.linalg.norm(V[k:n, k]). You should instead use V[k:n, k] = V[k:n, k] / np.linalg.norm(V[k:n, k]).

Computing Eigenvalues with QR Decomposition

We can leverage the QR decomposition to approximately compute the eigenvalues for a matrix $A \in \mathbb{R}^{n \times n}$. First, we let $A_{\text{temp}} = A$. Then, we iteratively update A_{temp} a total of m times, where the k-th iteration constitutes the following updates:

$$QR = A_{\text{temp}},$$

 $A_{\text{temp}} = RQ.$

Under certain conditions, A_{temp} will converge to an upper triangular matrix. Then, the eigenvalues of A are the diagonal elements of A_{temp} . The complete algorithm is presented below as Algorithm 2. Please use the qr function you implemented following Algorithm 1 for the QR decomposition step. Do not use NumPy's qr function.

Algorithm 2 Algorithm for Eigenvalue Computation

1: function COMPUTEEIGENVALUES(A, m)2: $A_{\text{temp}} \leftarrow A$ 3: for i = 0 to m - 1 do 4: Factor A_{temp} into $QR = A_{\text{temp}}$ 5: $A_{\text{temp}} \leftarrow RQ$

6: **return** diag (A_{temp})

Grading Breakdown

The grading breakdown for the assignment is as follows:

Problem 1	18%
Problem 2	18%
Problem 3	22%
Problem 4	11%
Programming	31%
Total	100%

Problem 0 is ungraded but must be submitted for your assignments this semester to be graded.

Handing In

You will turn in your final handin via Gradescope, which you should have been automatically added to through Canvas. If you have questions about using Gradescope or have not been added, please ask on Edstem. For this assignment, you should have written answers for Questions 1, 2, 3, and 4. Please submit the written portion to "Homework 1" on Gradescope.

For the programming, you should submit your code (warmup.py and eigenvalues.py) to "Homework 1 Code" on Gradescope.

Anonymous Grading

You need to be graded anonymously, so please do not write your name anywhere on your handin.

Obligatory Note on Academic Integrity

Plagiarism — don't do it.

As outlined in the Brown Academic Code, attempting to pass off another's work as your own can result in failing the assignment, failing this course, or even dismissal or expulsion from Brown. More than that, you will be missing out on the goal of your education, which is the cultivation of your own mind, thoughts, and abilities. Please review this course's collaboration policy and if you have any questions, please contact a member of the course staff.